

Name: Key

1. [4 parts, 1/2 point each] Determine the truth value of each of the following wffs in the interpretation where the domain is $\{0, 1, 2, 3, \dots\}$ and $P(x)$ is " x is even". (Remember that 0 is even.) Write the entire word "true" or the entire word "false".

(a) $\forall x [\exists y [P(y) \wedge y > x]]$

true

(b) $\exists x [\forall y [P(y) \rightarrow y > x]]$

false

(Note: even if $x=0$, it fails for $y=0$)

(c) $\exists x [\forall y [P(x) \rightarrow P(y)]]$

true;

(d) $\forall x [\exists y [P(x) \rightarrow P(y)]]$

true

2. [3 parts, 1 point each] Using the predicates shown, translate the following into wffs.

$J(x)$: " x is a judge"	$C(x)$: " x is a chemist"	$L(x)$: " x is a lawyer"
$W(x)$: " x is a woman"	$A(x, y)$: " x admires y "	

- (a) No woman is both a lawyer and a chemist.

$$\forall x [W(x) \rightarrow (L(x) \wedge C(x))']$$

Also OK: $(\exists x [W(x) \wedge L(x) \wedge C(x)])'$

- (b) Every judge admires a female lawyer.

$$\forall x [J(x) \rightarrow \exists y [W(y) \wedge L(y) \wedge A(x, y)]]$$

- (c) Some women admire only those who are lawyers or judges.

$$\exists x [W(x) \wedge \forall y [A(x, y) \rightarrow L(y) \vee J(y)]]$$

3. [2 points] Write the negation of the sentence "Some lawyers admire only judges". Your sentence should be as simple as possible.

Every lawyer admires someone who is not a judge.

Derivation Rule	Name/Abbreviation for Rule
$P \vee Q \iff Q \vee P$ $P \wedge Q \iff Q \wedge P$	Commutative—comm
$(P \vee Q) \vee R \iff P \vee (Q \vee R)$ $(P \wedge Q) \wedge R \iff P \wedge (Q \wedge R)$	Associative—ass
$(P \vee Q)' \iff P' \wedge Q'$ $(P \wedge Q)' \iff P' \vee Q'$	De Morgan's laws—De Morgan
$P \rightarrow Q \iff P' \vee Q$	Implication—imp
$P \iff (P')'$	Double negation—dn
$P \leftrightarrow Q \iff (P \rightarrow Q) \wedge (Q \rightarrow P)$	Defn of Equivalence—equ
$\left. \begin{array}{l} P \\ P \rightarrow Q \end{array} \right\} \implies Q$	Modus ponens—mp
$\left. \begin{array}{l} P \rightarrow Q \\ Q' \end{array} \right\} \implies P'$	Modus tollens—mt
$\left. \begin{array}{l} P \\ Q \end{array} \right\} \implies P \wedge Q$	Conjunction—con
$P \wedge Q \implies \left\{ \begin{array}{l} P \\ Q \end{array} \right.$	Simplification—sim
$P \implies P \vee Q$	Addition—add

4. [3 points] Using the given derivation rules and the 4 derivation rules involving quantifiers, give a proof sequence to show the following wff is valid.

$$\exists x [\forall y [Q(x, y)]] \rightarrow \forall y [\exists x [Q(x, y)]]$$

1. $\exists x [\forall y [Q(x, y)]]$ hyp
2. $\forall y [Q(a, y)]$ 1, ei
3. $Q(a, y)$ 2, ui
4. $\exists x [Q(x, y)]$ 3, eg
5. $\forall y [\exists x [Q(x, y)]]$ 4, ug