First-Fit chain partitions in partially ordered sets

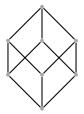
Kevin G. Milans (milans@math.wvu.edu)



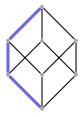
Michael C. Wigal

West Virginia University

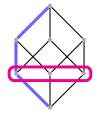
2018 International Workshop on Graph Theory
Ewha Womans University
Seoul, Korea
January 4, 2018



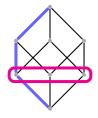
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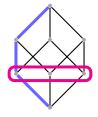
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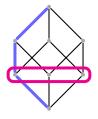
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The minimum size of a chain partition of P equals the width of P.

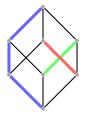


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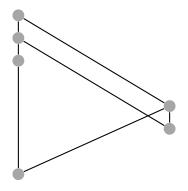
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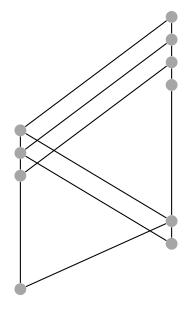
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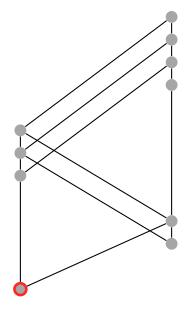
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- ▶ In terms of the width w, how many chains can First-Fit use?

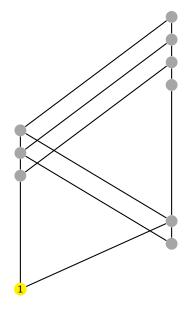


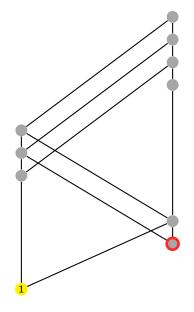


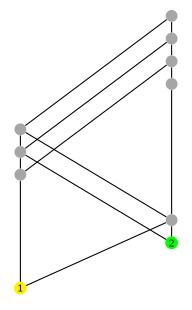


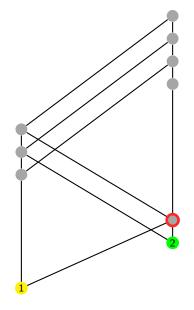


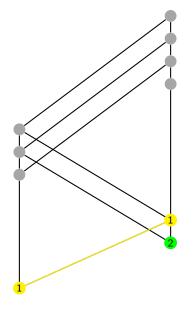


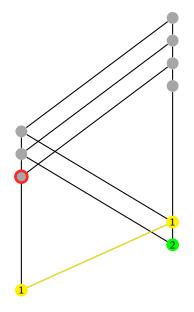


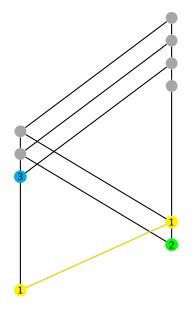


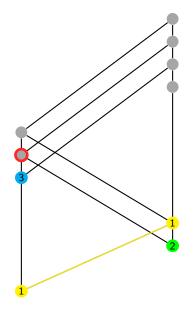


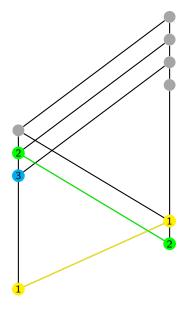


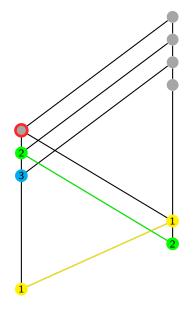


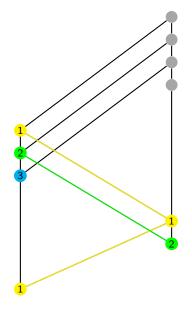


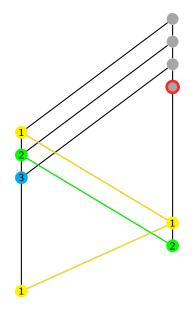


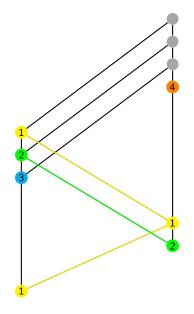


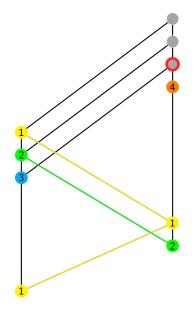


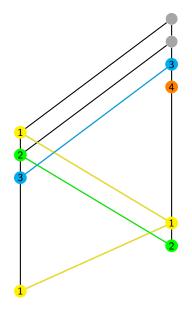


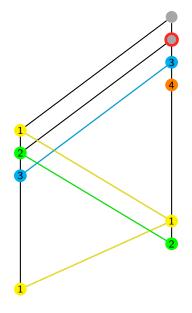


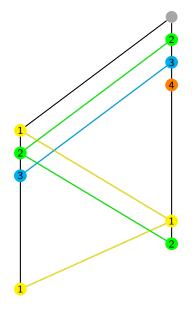


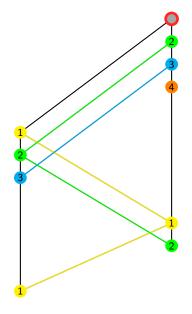


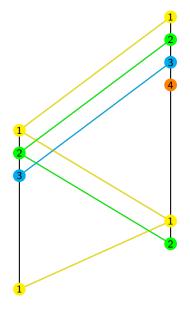












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Theorem (Bosek-Krawczyk-Matecki (2011))

If Q has width at most 2, then FF(w, Q) is bounded.



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- ▶ Let $P_1 + P_2$ be the poset with disjoint copies of P_1 and P_2 .



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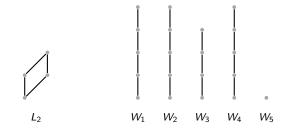
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Main Question

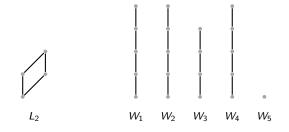
What properties of Q determine the behavior of FF(w, Q)?



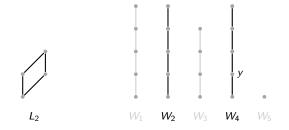
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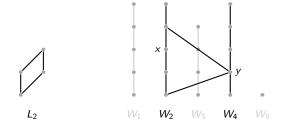
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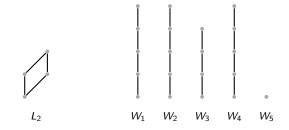
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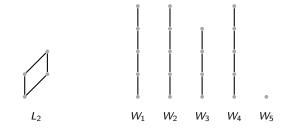
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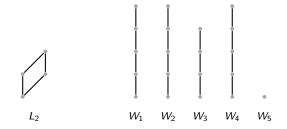
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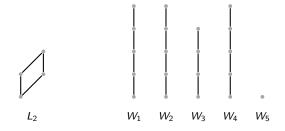
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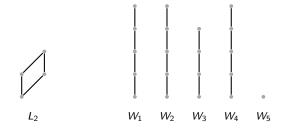
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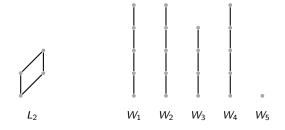
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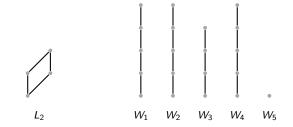
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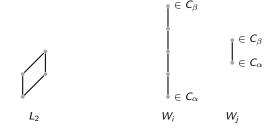
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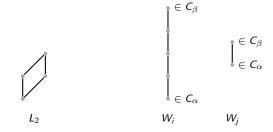
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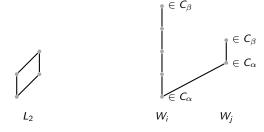
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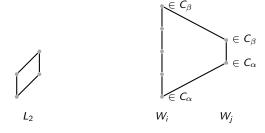
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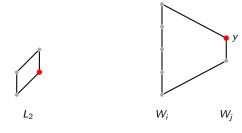
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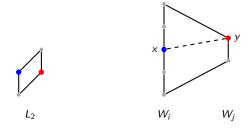
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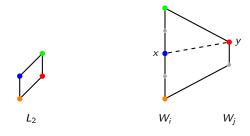
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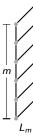


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- ▶ Similarly max W_i > min W_j .
- ▶ Let $y \in W_j$, and let $x \in W_i$ be incomparable to y.
- ▶ Note min $W_i < x, y < \max W_i$, completing a copy of L_2 .

Ladders



► The ladder poset, denoted L_m , has two chains $x_1 < ... < x_m$ and $y_1 < ... < y_m$ with $x_i \le y_j$ if and only if $i \le j$.



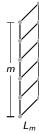
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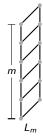
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- ▶ Bosek–Krawczyk (2015): if $FF(w, L_m)$ is polynomial in w and m, then there is a polynomial online chain partitioning algorithm.



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Theorem

$$\frac{1}{m-1} w^{\lg(m-1)} \le FF(w, L_m) \le w^{2.5 \lg 2w + 2 \lg m}$$

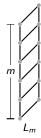


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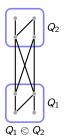


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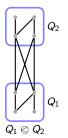


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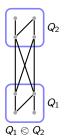




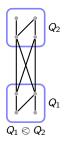
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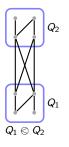


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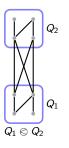


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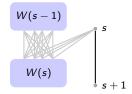
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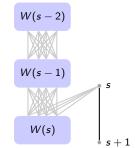
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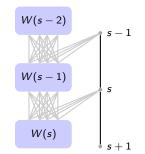
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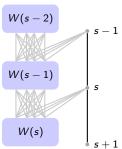
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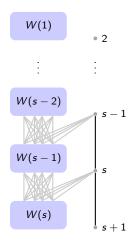


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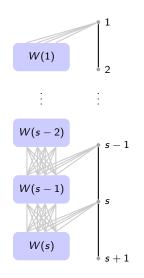


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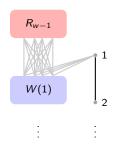


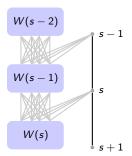


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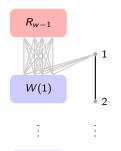


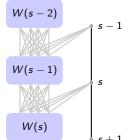
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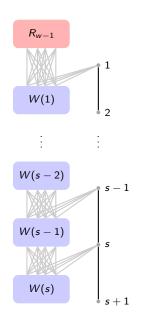


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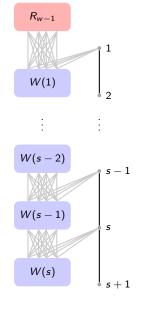




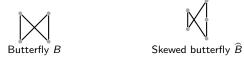
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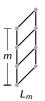
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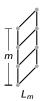
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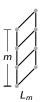
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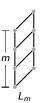
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