Turán and Ramsey Results for Boolean Algebras

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- Such a family of 2^d sets forms a copy of \mathcal{B}_d .
- A family is \mathcal{B}_d -free if it does not contain a copy of \mathcal{B}_d .

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- ► [Erdős–Kleitman 1971] For some constants c₁, c₂ and n sufficiently large

$$c_1 \cdot n^{-1/4} \cdot 2^n \leq b(n,2) \leq c_2 \cdot n^{-1/4} \cdot 2^n.$$

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$$b(n,d) \leq 50 \cdot n^{-\frac{1}{2^d}} \cdot 2^n.$$

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- For $d \ge 1$, we have $\binom{\alpha_d(n)}{2}/n = \alpha_{d-1}(n)$.

• Given x_0, x_1, \ldots, x_d with $x_0 \ge 0$ and $x_i \ge 1$ for $i \ge 1$,

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If $A \subseteq [0, n]$ and $|A| > \alpha_d(n)$, then A contains an affine d-cube.

• Using
$$\alpha_d(n) \le (4n)^{1-\frac{2}{2^d}} < 4n^{1-\frac{2}{2^d}}$$
, we obtain:

Corollary

If $A \subseteq [0, n]$ and $|A| \ge 4n^{1-\frac{2}{2^d}}$, then A contains an affine d-cube.

The Lubell Function



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- ► The Lubell function of *F*, denoted h_n(*F*), is E[X].
- ► Think of h_n(F) as a measure of the size of F, with 0 ≤ h_n(F) ≤ n + 1.

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where \mathcal{F}_S is the set of all $A \in \mathcal{F}$ that are disjoint from S with $A \cup S \in \mathcal{F}$.

Theorem If $\mathcal{F} \subseteq 2^{[n]}$ and $h_n(\mathcal{F}) > \alpha_d(n)$, then \mathcal{F} contains a copy of \mathcal{B}_d . Corollary (Szemerédi's Cube Lemma) If $A \subseteq [0, n]$ and $|A| > \alpha_d(n)$, then A contains an affine d-cube.

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• Note $\mathcal{F} \subseteq 2^{[n]}$ and $h_n(\mathcal{F}) = |\mathcal{A}| > \alpha_d(n)$.

▶ By the theorem: *F* contains a copy of *B_d* generated by disjoint sets *X*₀, *X*₁,..., *X_d*.

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- ▶ By the theorem: *F* contains a copy of *B_d* generated by disjoint sets *X*₀, *X*₁,..., *X_d*.
- ► Hence A contains an affine d-cube generated by x₀,..., x_d with x_i = |X_i|.

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- Open for $d \ge 2$.

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► By convexity: $\mathbf{E}[\binom{X}{2}] \ge \binom{\mathbf{E}[X]}{2} > \binom{\alpha_d(n)}{2} = n\alpha_{d-1}(n)$ ► Grouping pairs $(A, B) \in \mathcal{F} \times \mathcal{F}$ with $A \subsetneq B$ by B - A, with S = B - A: $\mathbf{E}[\binom{X}{2}] = \sum_{k=1}^{n} \frac{1}{\binom{n}{k}} \sum_{S \in \binom{[n]}{k}} h_{n-k}(\mathcal{F}_S),$ where \mathcal{F}_S is the family of all $A \in \mathcal{F}$ that

where \mathcal{F}_S is the family of all $A \in \mathcal{F}$ that are disjoint from S with $A \cup S \in \mathcal{F}$.

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If $\mathcal{F} \subseteq 2^{[n]}$ and $|\mathcal{F}| \ge 50n^{-1/2^d} \cdot 2^n$, then \mathcal{F} contains a copy of \mathcal{B}_d .

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